

# A partial list of mathematical symbols and how to read them

## Greek alphabet

A	$\alpha$	alpha	B	$\beta$	beta	$\Gamma$	$\gamma$	gamma	$\Delta$	$\delta$	delta	E	$\epsilon, \varepsilon$	epsilon
Z	$\zeta$	zeta	H	$\eta$	eta	$\Theta$	$\theta, \vartheta$	theta	I	$\iota$	iota	K	$\kappa$	kappa
$\Lambda$	$\lambda$	lambda	M	$\mu$	mu	N	$\nu$	nu	$\Xi$	$\xi$	xi	O	$o$	omicron
$\Pi$	$\pi, \varpi$	pi	P	$\rho, \varrho$	rho	$\Sigma$	$\sigma, \varsigma$	sigma	T	$\tau$	tau	$\Upsilon$	$\upsilon$	upsilon
$\Phi$	$\phi, \varphi$	phi	X	$\chi$	chi	$\Psi$	$\psi$	psi	$\Omega$	$\omega$	omega			

## Important sets

$\emptyset$	empty set	
$\mathbb{N}$	natural numbers	$\{0, 1, 2, \dots\}$
$\mathbb{N}^+$	positive integer numbers	$\{1, 2, \dots\}$
$\mathbb{Z}$	integer numbers	$\{\dots, -2, -1, 0, 1, 2, \dots\}$
$\mathbb{Q}$	rational numbers	$\{m/n : m \in \mathbb{Z}, n \in \mathbb{N}^+\}$
$\mathbb{R}$	real numbers	$(-\infty, +\infty)$
$\mathbb{R}^+$	positive real numbers	$(0, +\infty)$
$\mathbb{C}$	complex numbers	$\{x + iy : x, y \in \mathbb{R}\}$ ( $i$ is the imaginary unit, $i^2 = -1$ )

## Logical operators

$\forall$	for all, universal quantifier	$\forall n \in \mathbb{N}, n \geq 0$
$\exists$	exists, there is, existential quantifier	$\exists n \in \mathbb{N}, n \geq 7$
$\exists!$	there is exactly one	$\exists! n \in \mathbb{N}, n < 1$
$\wedge$	and	$(3 > 2) \wedge (2 > 1)$
	...over an index set	$\bigwedge_{i \in \mathbb{N}} B_i = B_0 \wedge B_1 \wedge B_2 \wedge \dots$
$\vee$	or	$(2 > 3) \vee (2 > 1)$
	...over an index set	$\bigvee_{i \in \mathbb{N}} B_i = B_0 \vee B_1 \vee B_2 \vee \dots$
$\Rightarrow$	implication, if-then	$\forall a, b \in \mathbb{R}, (a = b) \Rightarrow (a \geq b)$
$\Leftrightarrow$	biimplication, if-and-only-if	$\forall a, b \in \mathbb{R}, (a = b) \Leftrightarrow (b = a)$
$\neg$	negation, not	$\neg(2 > 3)$
	alternative notations for negation	$\overline{(2 > 3)}, 2 \not> 3$

## Arithmetic operators

$  $	absolute value	$ -7  =  7  = 7$
$\sum$	summation	$\sum_{i \in \mathbb{N}^+} 2^{-i} = 1$
$\prod$	product	$\prod_{i=1}^n i = n!$
$!$	factorial	$7! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 \cdot 6 \cdot 7 = 5040$
$\binom{n}{m}$	$n$ choose $m$ , combinatorial number	$\binom{n}{m} = \frac{n!}{(n-m)!m!}$
mod	modulo, remainder	$7 \bmod 3 = 1, -8 \bmod 5 = 2$
div	integer quotient	$7 \operatorname{div} 3 = 2, -8 \operatorname{div} 5 = -2$

## Set operators

$\in$	in, membership	$a \in \{a, b, c\}$
$\cup$	union	$\{a, b, c\} \cup \{a, d\} = \{a, b, c, d\}$
	... over an index set	$\bigcup_{i \in \mathbb{N}} S_i = S_0 \cup S_1 \cup S_2 \cup \dots$
$\cap$	intersection	$\{a, b, c\} \cap \{a, d\} = \{a\}$
	... over an index set	$\bigcap_{i \in \mathbb{N}} S_i = S_0 \cap S_1 \cap S_2 \cap \dots$
$\setminus$	difference	$\{a, b, c\} \setminus \{a, d\} = \{b, c\}$
$\supset$	strict superset	$\mathbb{Z} \supset \mathbb{N}$
$\supseteq$	superset	$\mathbb{N} \supseteq \mathbb{N}$
$\subset$	strict subset	$\mathbb{N} \subset \mathbb{Z}$
$\subseteq$	subset	$\mathbb{N} \subseteq \mathbb{N}$
$2^A$	power set of $A$	if $A = \{a, b, c\}$ , then $2^A = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, A\}$

## String, grammar, and formal language notation

$\lambda$	empty string (at times, $\epsilon$ is used instead of $\lambda$ )	$\lambda a = a$
$*$	Kleene star, zero or more occurrences	$a^* = \{\epsilon, a, aa, aaa, \dots\}$
$+$	one or more occurrences	$a^+ = \{a, aa, aaa, \dots\}$
$ w $	length of string $w$	$ abc  = 3,  a^n  = n,  \epsilon  = 0$
$ w _a$	number of occurrences of $a$ in string $w$	$ aab _a = 2,  aab _b = 1,  aab _d = 0$
$A \rightarrow x$	$A$ goes to $x$ (grammar production)	
$A \Rightarrow x$	$A$ derives $x$	
$A \xRightarrow{*} x$	$A$ derives $x$ in a number of steps	
$A \xrightarrow[G]{\Rightarrow} x$	$A$ derives $x$ according to $G$	
$A \xrightarrow[G]{\xRightarrow{*}} x$	$A$ derives $x$ according to $G$ in a number of steps	
$(q, aa) \vdash (p, a)$	$(q, aa)$ yields $(p, a)$ in one step	
$(q, aa) \vdash^* (p, a)$	$(q, aa)$ yields $(p, a)$ in a number of steps	
$(q, aa) \vdash_M (p, a)$	$(q, aa)$ yields $(p, a)$ in one step according to $M$	
$(q, aa) \vdash_M^* (p, a)$	$(q, aa)$ yields $(p, a)$ in a number of steps according to $M$	
$M \searrow w$	the Turing machine $M$ halts on string $w$	
$M \nearrow w$	the Turing machine does not $M$ halt on string $w$	

## And remember...

$0! = 1$
$\forall n \in \mathbb{Z}, \forall m \in \mathbb{N}, m > 0 \Rightarrow n = (n \operatorname{div} m)m + (n \operatorname{mod} m)$
$\bigcup_{i \in \emptyset} S_i = \emptyset$
$\sum_{i \in \emptyset} n_i = 0$
$\prod_{i \in \emptyset} n_i = 1$